Next-Generation Debuggers

For Reverse Engineering

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The ERESI project

- Started in 2001 with the ELF shell
- Developed at LSE (EPITA security laboratory)
- Contains more than 10 components
- Featured in 2 articles in Phrack Magazine:
 - The Cerberus ELF Interface (2003)
 - Embedded ELF Debugging (2005)

Limitations of existing UNIX debugging framework

- GDB: Use OS-level debugging API (ptrace) -> does not work if ptrace is disabled or absent
- Very sensitive to variation of the environment (ex:
 - ET_DYN linking of hardened gentoo)
 - Strace / Ltrace : use ptrace as well. Very few interaction (command-line parameters)

Limitations of existing frameworks

None of these frameworks rely on a real reverse engineering language

The ERESI team

- Started with a single person in 2001 (The ELF shell crew). Remained as it during 3 years.
- Another person developed libasm (disassembling library) since 2002
- A third person developed libdump (the network accessibility library) in 2004-2005
- Since mid-2006 : community project (6 persons)

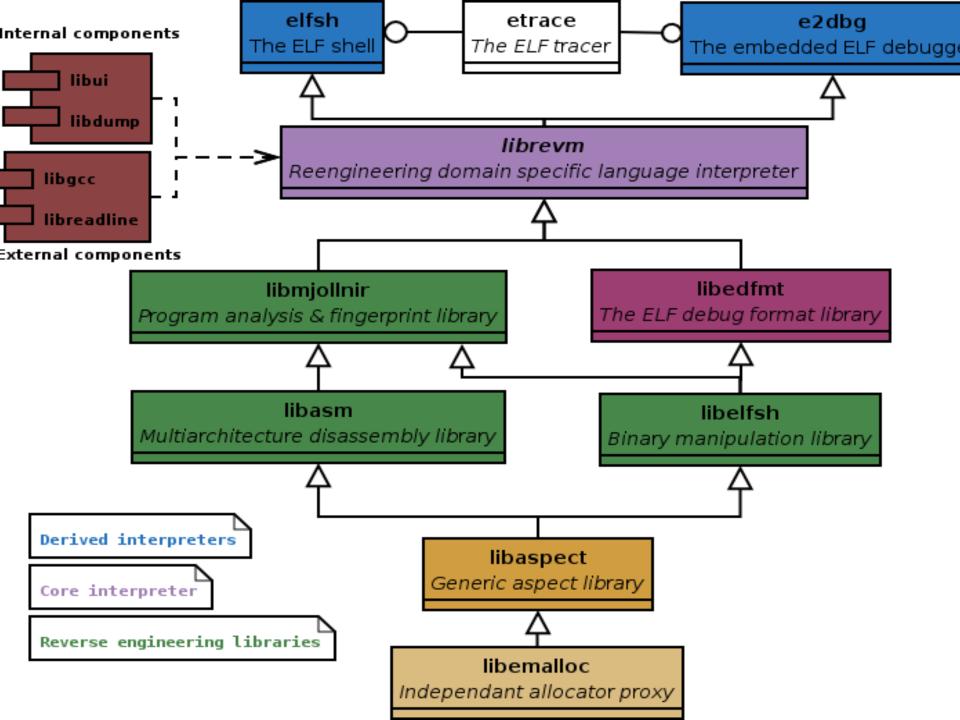
The modern ERESI project

- elfsh (and libelfsh): The ELF shell
- e2dbg (and libe2dbg): The Embedded ELF
- debugger
- etrace : The Embedded tracer
- librevm : the language interpreter
- libmjollnir : fingerprinting & graphs library
 - libaspect: Aspect oriented library

The modern ERESI project (cont)

- libasm : the disassembling library
- libedfmt : the ERESI debug format library
- liballocproxy : allocation proxying library
- libui : The user interface (readline-based)





ERESI contributions (1)

Can debug hardened systems (does not need ptrace): PaX/grsec compatible

Very effective analysis: improve the performance of fuzzing and heavy-weight debugging (no context switching between the debugger and the debuggee: the dbgvm resides in the debuggee)

ERESI contributions (2)

- A reflective framework : possibility to change part of it in runtime without recompilation
- The first real reverse engineering language!
 - hash tables
 - regular expressions
 - loops, conditionals, variables
 - The complete ELF format objects accessible from the language

The ERESI language: example 1

load /usr/bin/ssh

```
set $entnbr 1.sht[.dynsym].size
div $entnbr 1.sht[.dynsym].entsize
print Third loop until $entnbr :
foreach $idx of 0 until $entnbr
    print Symbol $idx is 1.dynsym[$idx].name
forend
```

unload /usr/bin/ssh

The ERESI language: example 2

```
add $hash[hname] Intel
add $hash[hname] Alpha
add $hash[hname] Sparc32
add $hash[hname] Mips
add $hash[hname] Sparc64
add $hash[hname] AMD
add $hash[hname] Pa-risc
foreach $elem of hname matching Sparc
    print Regex Matched $elem
endfor
```



The ERESI language: example 3

```
type archtypes = elm:string[55]
inform archtypes elfsh_arch_type
type archaddr = elm:long[55]
inform archaddr elfsh_arch_type
print Now print Strings
print 107.archtypes[elfsh_arch_type].elm[0]
print 107.archtypes[elfsh_arch_type].elm[1]
print Now print addresses
print 107.archaddr[elfsh_arch_type].elm[0]
print 107.archaddr[elfsh_arch_type].elm[1]
```

Note: elfsh_arch_type is the symbol of an array in elfsh. In that example, elfsh itself is analysed

e2dbg: the Embedded ELF debugger

- Does not use ptrace. Does not have to use any OS level debug API. Evades PaX and grsecurity.
- Proof of concept developed on Linux / x86.
- Scriptable using the ERESI language
- Support debugging of multithreads
- No need of ANY kernel level code (can execute in hostile environment)



- + Unintrusive heap
- + analysis code
- + aspect library
- + debug format handling

Client-side debugger

- Target abstraction
- Communication abstraction
- Interface abstraction

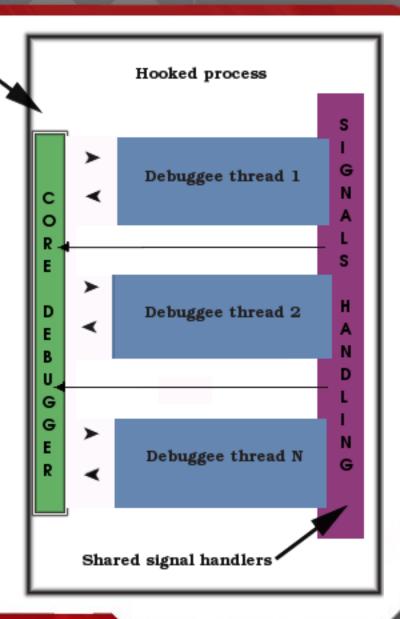


FIFO INET (...)





- Interprocess communication
 - ➤ Intraprocess communication



e2dbg: features

- Classical features:
 - breakpoints (using processor opcode or function redirection)
 - stepping (using sigaction() syscall)
- Allocation proxying
 - keep stack and heap unintrusiveness
- Support for multithreading

Allocation proxying

We manage two different heap allocator in a single process:

```
int hook_malloc(int sz)
{
    if (debugger)
      return (aproxy_malloc(sz));
    return (orig_malloc(sz))
}
```

Handling of debug format in ERESI

Debugging format

- Describe each element of a program
 - Give names and position of:
 - Variables
 - Functions
 - Files
 - •
 - Store program types dependences between them

Debugging format - issues

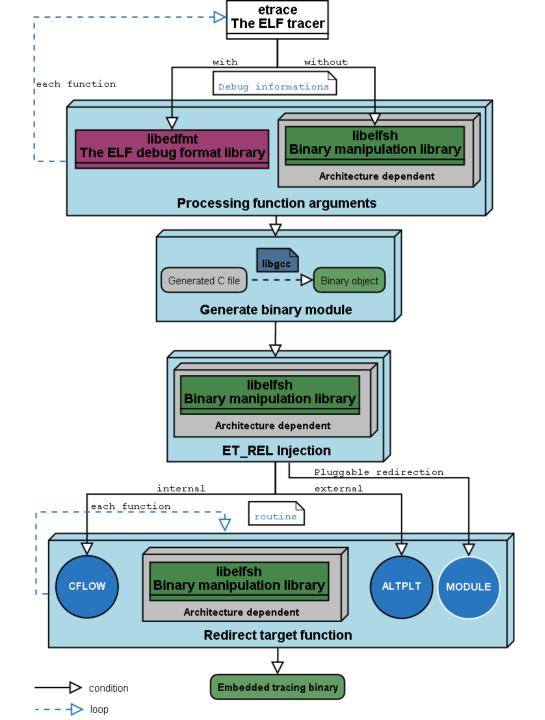
- Distinction of debugging format
 - stabs, dwarf, stabs+, dwarf2, gdb, vms ...
 - Different ways to parse, read, store ...
 - For example with stabs and dwarf2
 - Stabs does not contain any position reference
 - You store the whole parsing tree
 - Dwarf2 use read pattern apply directly on data
 - You cannot store everything (too big)
 - **—** ...

Uniform debugging format

- Parsing
 - So we can read the debugging format
- Transforming
 - We transform it on a uniform representation
 - Keep only useful information
- Cleaning
 - We keep only uniform debugging format
 - New debugging format
 - We change only backward part
- Register types on ERESI type engine

Embedded ELF tracer

- Tracer using ELFsh framework
- Tracing internal and external calls
- Dynamic and supports multiple architecture
 - It does not use statically stored function prototypes
 - Use gcc to reduce architecture dependence
- Work with and without debugging format
- Recognize string, pointers and value



Embedded ELF tracer

- Trace backend
 - Analyze target function
 - Create proxy functions
- Embedded tracer
 - Inject proxy functions in the binary
 - Redirect calls into our proxy functions
 - Create a new binary
 - Automatic using the ELF tracer

Embedded ELF trace - script

#!/usr/local/bin/elfsh32

load ./test

traces add main

traces create critical

traces add crypt critical

traces create mainnodes

traces add firstfunc mainnodes

traces add testcrypt mainnodes

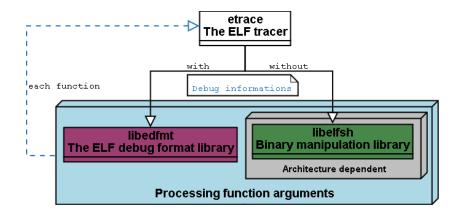
save test2

Embedded ELF trace - output

```
+ main()
BEFORE!
 + firstfunc(int num: 0x3, *char value: *0x8048648 "this is the text")
arguments: num = 3 / value = this is the text (0x8048648)
   + testcrypt()
    + crypt(*0x804860d "password", *0x8048608 "salt")
    - crypt = b7f74120
   - testcrypt = b7f74120
 - firstfunc = 1
OK!
AFTER
- main = 0
```

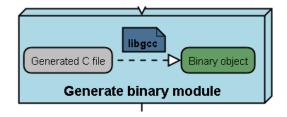
Etrace - Processing function arguments

- With debugging information
 - Extract arguments information
 - size
 - names
 - type names
 - ...
- With architecture dependent argument counting
 - Backward analysis
 - Forward analysis



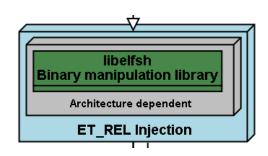
Etrace - Generate binary module

- Generate a .c file
 - Call tree (padding)
 - Dynamic check pointers, strings or value
- Benefits
 - Architecture independent
 - New feature implementation
 - Less bugs
 - Use ELFsh framework



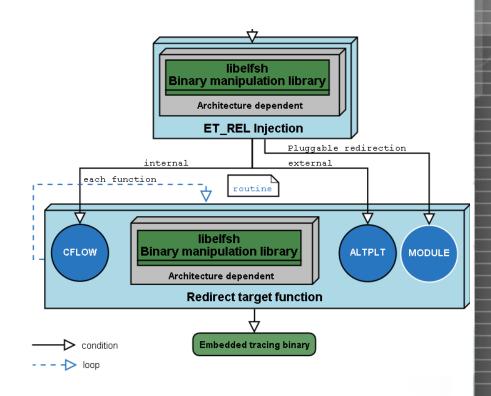
Libelfsh - ET_REL injection

- ET_REL injection principle
 - Add a binary module directly on target binary
- Merge symbols and sections list
- Section injection
 - Code sections
 - Injected before .interp
 - Data sections
 - Injected after .bss
 - Relocation in two steps



Libelfsh - Redirect target function

- Internal function
 - CFLOW technique
- External function
 - ALTPLT technique
- Custom redirection
 - Vector benefit
 - Your own redirection mechanism



Program analysis

A Graph Analyzer

- Graph analyzers
 - Identify blocks and functions
 - Identify links (calls and jumps)
 - Build a graph with this info
- Control Flow Graphs (CFGs)
 - Inter-blocks CFGs vs. Interprocedural CFGs
 - Main instrument to Control Flow analysis

A Graph Analyzer

- **Control Flow Analysis**
 - Essential to some kinds of further analysis and to optimization
 - Gives information about properties such as
 - Reachability
 - Dominance
 - •

A Graph Analyzer – Libasm

Libasm

- Lowest layer of this application
- Multi-architecture disassembling library
 - Intel IA-32
 - SPARC V9
 - In the near future, MIPS
- Unified type system

A Graph Analyzer – Libasm

| Trme | Description |
|------------|---|
| Type | Description |
| IMPBRANCH | Imperative branch (jump) |
| CONDBRANCH | Conditional branch |
| CALLPROC | Call to a procedure |
| RETPROC | Return from a procedure |
| ARITH | Arithmetic or logic operations |
| LOAD | Memory data load |
| STORE | Memory data store |
| ARCH | Architecture-dependent instruction |
| FLAG | Flag-modifier instruction |
| INT | Interrupt or call-gate instruction |
| ASSIGN | Assignment instruction |
| TEST | Comparison or test instruction |
| NONE | Instruction that doesn't fit any of the above |

A Graph Analyzer – Libasm

- The unified instruction type system
 - Works with non-mutually exclusive types
 - Provides means to "blindly" analyze an instruction
 - Eg. Control Flow analysis!

A Graph Analyzer - Libasm

- Libasm vectors
 - Storage of pointers to opcode handling functions
 - 4 dimensions: 1 for machine info, 3 for opcode info
 - Runtime dumping and replacing of vectors
 - Built-in language constructs
 - Easy-made opcode tracer!

A Graph Analyzer – libmjollnir

- Libmjollnir
 - Upper-layer component
 - Code fingerprinting and program analysis
- CFG construction
 - Libmjollnir treats both: blocks and functions
 - Separate representations (structures)

A Graph Analyzer – libmjollnir

Containers

- Generic structures to encapsulate blocks and functions
- Have linking (input and output links) information
- Have a pointer to data and type information to interpret this data accordingly

A Graph Analyzer – libmjollnir

Containers

- Allow for more abstract graph analysis (analyzing a graph of containers)
- In the future, may also store data nodes (Data Flow analysis)
- Also for the future, containers of containers
 - Even higher abstraction of links and relationships

Questions?

Thank you for your attention

If you are interested in joining us, come to talk after the conference.